

SUPPLEMENTARY PROBLEMSCHAPTER 1 - LANGUAGES & REG. EXPR.

E1. Let $A = \{a\}$ and $B = \{b\}$. Describe the following sets

(a) A^*

(e) $(A \cap B)^*$

(b) $A^* \cdot B^*$

(f) $(A \cup B)^* \cdot AB$

(c) $(A \cdot B)^*$

(g) $(A^* \cup AB)^* \cdot A \cup B^*$

(d) $(A \cup B)^*$

(h) $(A^* \cup AB)^* - (BAB)^*$

E2 Find the simplest expressions for the following

(a) $\lambda \cdot \emptyset^*$

(h) $\{\emptyset, \{\emptyset\}\} - \{\emptyset\}$

(b) $\lambda^* \cdot \emptyset^*$

(i) $\{\emptyset\} \cap \{\emptyset, \{\emptyset\}\}$

(c) $A^* \cup \emptyset^*$

(j) $\{\emptyset\} \cap \emptyset$

(d) $(\lambda \cup A)^*$

(k) $\lambda - \emptyset^*$

(e) $(\emptyset \cup A)^*$

(l) $\lambda - \{\emptyset^*\}$

(f) $(\lambda^* \cdot \emptyset^*)^*$

(m) $\lambda^* - \emptyset^*$

(g) $\{\emptyset, \{\emptyset\}\} - \emptyset$

(n) $\lambda \cdot \emptyset$

E3 Find regular expressions for the following languages

(a) $\{\varphi \in \{0,1\}^* : 101 \text{ is a substring of } \varphi\}$

(b) $\{\varphi \in \{0,1\}^* : \text{each } 0 \text{ in } \varphi \text{ is immediately followed by at least two } 1\text{'s}\}$

(c) $\{\varphi \in \{0,1\}^* : \text{each } 1 \text{ in } \varphi \text{ is immediately followed by the string } 10\}$

(d) $\{\varphi \in \{0,1\}^* : 101 \text{ is not a substring of } \varphi\}$

E4. Say which of the following equations are always true & provide proofs to justify your answers. Also say which of the equations are not always true and provide counterexamples to justify your answers. A, B & C are arbitrary languages.

$$(a) A(BA)^* = (AB)^*A$$

$$(b) (AB)^* = A^*B^*$$

$$(c) A \cdot (B^* \cap C^*) = (A \cdot B^*) \cap (A \cdot C^*)$$

$$(d) A^* \cdot B = B \cup A^* \cdot AB$$

$$(e) (A \cap B)^* = A^* \cap B^*$$

$$(f) A^* \cdot (B \cap C)^* = (AB \cup AC)^*$$

$$(g) (A^*B^*)^* = (A \cup B)^*$$

$$(h) A^*(B \cup C) = A^*B \cup A^*C$$

E5. Let S be a set of strings of letters from the alphabet V . We say that S is commutative if for any α & β in S

$$\alpha \cdot \beta = \beta \cdot \alpha$$

(a) Prove that if $S \subseteq (\underline{w})^*$, then S is commutative. Here w = a fixed string and $(\underline{w})^* = \{w^n : n \geq 0\}$

(b) Prove that if S is commutative then we can find a string w such that $S \subseteq (\underline{w})^*$

CHAPTER 4 - REGULAR LANGUAGES

E6. (a) Prove that the language $\{a^k b^{3k} : k \geq 1\}$ is non-regular

(b) Prove that the language $\{\varphi c \varphi : \varphi \in \{0,1\}^*\}$ is non-regular.

E7. A prime number is any positive integer $p > 1$ which has only 1 & p as its divisors.
Let $L = \{a^p : p \text{ is a prime number}\}$.
Prove that L is a non-regular language.

E8. Let X and Y be regular languages based on the alphabet V . Determine which of the following languages are always regular and which of them are not always regular.

(a) $\{w : w \in X \text{ and } w^R \in Y\}$

(b) $\{w : w \in X \text{ and } w^R \notin Y\}$

(c) $\{w : w \in X \text{ and } w^R = w\}$

(d) $\{\varphi : \varphi^R \in X \text{ or } \varphi \notin Y\}$

(e) $\{\varphi : \varphi^R \notin X \text{ and } \varphi \in X.Y\}$

E9. Determine which of the following languages are regular and which are non-regular

(a) $\{a^k c a^l c a^m : m = k+l\}$

(b) $\{a^k c a^l c a^m : m \equiv k+l \pmod{3}\}$

(c) $\{a_1 \dots a_{2n} \in \{0,1\}^* : a_1 \dots a_n = a_{n+1} \dots a_{2n}\}$

(d) $\{w \in \{0,1\}^* : N_0(w) - N_1(w) \text{ is an even positive integer}\}$

Here $N_0(w)$ = number of 0's in w .

and $N_1(w)$ = number of 1's in w .

E10. (a) Let $L = \{P 0^n Q 1^n R : n \geq 1\}$ where P, Q & R are non-empty languages based on $\{0,1\}$. Is there a choice of P, Q and R for which L is regular

(b) Find infinite sets $L_1 \subsetneq L_2 \subsetneq L_3$ such that L_1 & L_3 are non-regular and L_2 is regular.

(c) Find infinite sets $L_1 \subsetneq L_2 \subsetneq L_3$ such that L_1 & L_3 are regular and L_2 is non-reg.

E11. Suppose A & B are non-regular sets. Determine if it is possible for any of the following sets to be regular.

(a) $A - B$

(b) $A \cup B$

(c) $A \cdot B$

CHAPTER 6 - RECURSIVE FUNCTIONS

E12. Define $\text{sign}(n) = \begin{cases} 0 & \text{if } n=0 \\ 1 & \text{if } n>1 \end{cases}$ and

$$\overline{\text{sign}}(n) = \begin{cases} 1 & \text{if } n=0 \\ 0 & \text{if } n>1. \end{cases}$$

Show that sign & $\overline{\text{sign}}$ are primitive recursive

E13. Show that each of the following functions, that are defined below, are primitive recursive.

(i) $\text{EXP}(x, y) = x^y$

(ii) $\text{ABS}(x, y) = |x - y|$

(iii) $\text{ZER}(x) = \begin{cases} 1 & \text{if } x=0 \\ 0 & \text{if } x>0 \end{cases}$

(iv) $\text{MIN}(x, y) = \text{SMALLER OF } x \text{ \& } y$

(v) $\text{MAX}(x, y) = \text{LARGER OF } x \text{ \& } y$

(vi) $\text{REM}(x, y) = \text{Remainder after dividing } y \text{ by } x$

(vii) $\text{QUO}(x, y) = \text{quotient obtained by dividing } y \text{ by } x$

(viii) $\text{EQU}(x, y) = \begin{cases} 1 & \text{if } x=y \\ 0 & \text{if } x \neq y. \end{cases}$

E14. Define $\text{ls}(x, y) = \begin{cases} 1 & \text{if } x < y \\ 0 & \text{if } x \geq y \end{cases}$ and

$$\text{gr}(x, y) = \begin{cases} 1 & \text{if } x > y \\ 0 & \text{if } x \leq y. \end{cases}$$

Show that ls and gr are primitive recursive functions.